METHOD AND APPARATUS FOR SYNCHRONIZING DATA TRANSFER

Related Application

     This application is a Non-Provisional Application based on
United States Provisional Application No. 60/244391, the disclosure
of which is incorporated herein by reference.

Field of the Invention

     The invention pertains to digital communication between two
circuit components. The invention is particularly adapted for use
in connection with the Advanced Communication Riser Integrated
10 Packet Bus communication specification.

Background of the Invention

     In many digital communication protocols, and particularly time
division multiplex (TDM) protocols, digital data is transmitted in
15 frames, where a frame comprises a plurality of time slots, each

time slot corresponding to a transmitted digital symbol (e.g., a

bit). As is well known in the art, in a frame, some of the time

slots may be dedicated to communication protocol housekeeping

information while the remaining time slots are dedicated to actual

5   data communicated between two digital devices. The housekeeping

data is usually transmitted at the beginning or the end of the

frame. Also as is well known to those of skill in the art, in TDM

systems, one or more time slots in a frame are dedicated to a

particular communication session between two particular digital

10 devices, while other time slots are dedicated to other

communication sessions. Accordingly, multiple communication

sessions can be carried on simultaneously.

Many different communication protocols are in use for TDM and

other frame-based digital communication paradigms. However, most,

15 if not all, include at least three signal lines, namely, a clock

line, a frame synchronization line and a data line. All

communication timing is keyed off of the clock, which typically

runs at the communication data bit rate of the system or a multiple

thereof. Obviously, the actual data being exchanged between two

20 devices is transmitted on the data line (if serial) or lines (if

parallel). The frame synchronization (or frame sync) line carries

a signal that delimits the beginning of each frame. Commonly, the

beginning of the frame is delimited by a rising or falling edge on

the frame sync line. All of the transmitting and receiving devices

25 synchronize themselves to the frame sync signal so that they will

be able to transmit and receive the symbols within the appropriate
time slots.

   In accordance with many digital communication protocols, the
clock and or frame sync signal is always generated by the
5   transmitting device. Other protocols allow or even require one
particular device to generate the clock and/or the frame sync
signal for another device (or devices). One such protocol is the
Integrated Packet Bus (IPB) specification promulgated by the
Advanced Communication Riser (ACR) Special Interest Group (SIG).
10   The ACR IPB specification is a protocol particularly designed for
communication between a CPU in a personal computer and peripheral
devices, particularly communications peripherals, such as home
networking hubs, ethernet hubs, modems, DSL modems, cable modems,
etc. The ACR IPB specification speaks in terms of a controller and
15   one or more targets. As is well known in the art of general
purpose, personal computer design, the trend has been to push out
all handling of communications between the CPU and other devices,
including peripherals to a dedicated chip, termed the controller.
In accordance with the IPB specification, part of the functionality
20   of the controller would be to carry out communications with
appropriate target devices, such as the aforementioned modems and
networking hubs using the ACR IPB protocol. Typically, the
controller will provide the functionality for communicating with
other devices also. Under the ACR IPB specification, the bus
25   between the controller and the one or more target devices comprises
six lines, namely, a received direction clock (RDCLK), a transmit

direction clock (TDCLK), a receive direction frame synchronization
signal (RDFRAME), a transmit direction frame synchronization signal
(TDFRAME), a transmit direction data line (DATAOUT) for data
transmitted from the controller to the target device, and a receive
5 direction data line (DATAIN). In ACR IPB, the transmit direction
is the direction from the controller to the target device and the
receive direction is the direction from the target device to the
controller. Under the ACR IPB protocol, the TDCLK and TDFRAME
signals are used for timing in the transmit direction are generated
10 at the controller. In addition, the RDFRAME signal is generated at
the controller. Under the ACR IPB protocol, the RDCLK signal can
be generated by either the controller or the transmitting target
device.

In digital communication protocols in which the clock or frame
15 sync signal for given set of transmission data is generated by a
device other than the transmitting device, a delay skew issue
arises.

For instance, consider the timing diagrams of Figures 1A and
1b. Figure 1A illustrates an exemplary timing situation for
20 transmit direction data in accordance with the ACR IPB protocol.
In accordance with the protocol, the first bit of data, B0 in a
frame begins being transmitted by the controller at time T_1 at the
beginning of the second clock cycle after the leading edge of the
frame synchronization signal. In this example, the leading edge of
25 the frame sync signal occurs at time T_0 and, for sake of
illustrating a worst case scenario, the first clock edge on TDCLK

is essentially simultaneous therewith (but late enough to be counted as the first clock edge after the frame sync). Accordingly, one full clock cycle later, at time T_1, the controller begins to transmit the first bit of data. Thus, that data bit
5 becomes fully valid at time $T_1+\delta$. At the falling edge of that clock cycle, at time T_2 (note that we are assuming a 50% duty cycle for TDCLK), the receiving target device samples the DATAOUT line to read the bit value. The process then continues for each clock cycle until the end of the frame.

10   Thus, for example, at T_3 (the third rising edge of TDCLK after the beginning of the frame), the controller begins transmitting the second bit. That second bit becomes fully valid at time $T_3+\delta$ and then is sampled by the receiving target device at time T_4 (the falling edge half-way through that clock cycle). Since the
15 TDFRAME, TDCLK and DATAOUT signals are all sourced from the same location (i.e., the controller), there are no delay skew issues. Particularly, even though there is some finite time lapse between the transmission of TDCLK, TDFRAME and DATAOUT from the controller and their receipt at the receiving target device, the delay is
20 essentially equal for all three lines.

However, in the receive direction, if the RDCLK and RDFRAME signal are sourced from the controller, while the DATAIN signal is sourced from the transmitting target device, delay skew becomes an issue. For instance, let us consider the example of Figure 1B.
25 The timing in the example of Figure 1B differs from that of Figure 1A in two respects. First, the clock and frame signals are

sourced from the controller, whereas the data is sourced from the
transmitting target device. Secondly, the first bit of data goes
out on the next clock edge (i.e., half of a clock cycle) after the
first clock edge after the frame sync signal, instead of a full

5 clock cycle after the first clock edge after the frame sync signal.

This situation presents two potential timing problems. First,
if the clock rate is very fast, the transmitting target device
simply may not be able to generate the first bit of the frame
within the allotted time. Particularly, Figure 1B illustrates

10 something close to a worse case scenario in which the first clock
edge after the frame sync signal occurs essentially simultaneously
with the frame sync signal at time T_0, but late enough to be
counted as the first clock edge after the frame sync. Then the
transmitting target device has only until T_1, i.e. half a clock

15 cycle later, to output the first bit of data. Some delay past T_1
is acceptable, because that data really only needs to be fully
valid by time T_2 when the controller will read the data.
Accordingly, the target device has slightly less than one full
clock cycle after the frame sync to output fully valid data. At

20 very high clock speeds, this may be very difficult to achieve. For
instance, the ACR IPB specification specifies a 40 MHz data rate,
which could make it very difficult for a transmitting target device
to place the data in the first bit position of a frame in time.

A second timing issue arises due to delay skew since the clock

25 and frame signals are sourced from the controller, but the data is
sourced from the target device. Particularly, as previously noted,

the transmitting device will begin transmitting a bit at the

leading edge of a clock cycle and that data must be fully valid at

the transmitting device at the midpoint edge of that clock cycle.

However, there is some finite delay between the time the controller

5   generates the leading edge of that clock cycle and the time that it

is received at the target device. Further, there is a delay

between the time the transmitting target device generates the data

and the time it is received at the receiving controller. The sum

of these two delays is called round trip delay. The round trip

10 delay includes not only the path delay but printed circuit board

delay and, possibly, a connector delay as well as internal device

delays. If the round trip delay is greater than half a clock cycle

(assuming a 50% duty cycle), the data cannot arrive at the

controller by the time the controller samples for that bit.

15   Thus, referring to Figure 1B, the controller experiences the

first clock edge after the frame sync at time T_0, and expects the

target device to begin transmitting the data at time T_1, and will

sample for the data at time T_2, just as in Figure 1A. However, the

target device experiences the first clock edge after the frame sync

20 at time $T_0+\Delta$ and begins transmitting the bit at time $T_1+\Delta$. Even

further, the controller actually receives the beginning of the

transmission of the data at time $T_1+2\Delta$. Of course, we also must

add in the delay, δ, between the transmitting device commencing

transmission of that bit and the bit value becoming fully valid.

25 Accordingly, the data is not fully valid at the receiving

controller until time $T_1+2\Delta+\delta$. If $2\Delta+\delta$ is greater than a half

clock cycle (as shown in Figure 1B), the data will not reach the
controller after time T_2 when the controller looks for the data.

Summary of the Invention

5   The invention is a method and apparatus for ensuring
synchronization for digital communication between a transmitting
and a receiving device, particularly when the clock and/or frame
synchronization signal is sourced from a different location than
the transmit data. In accordance with a first aspect of the
10 invention, the transmitting device employs a look ahead algorithm
to begin transmitting data before actual receipt of the frame
synchronization signal. Particularly, after the first frame
synchronization signal is received, the transmitting device has an
expectation of when the next (and each subsequent) frame
15 synchronization signal will be received and begins transmitting the
data in anticipation of receipt of the frame synchronization
signal.

        In accordance with a second aspect of the invention, at the
initialization of a communication session between two particular
20 devices, e.g., a controller and a target device, a test
communication session is conducted in which the receiving target
device receives an instruction from the controller to transmit a
test pattern (i.e., a series of consecutive bits of known pattern)
at the beginning of the next frame. The controller then looks for
25 receipt of that test pattern using the designated time slots in
accordance with the protocol (assuming no delay skew). If it does

not receive that test pattern, then it does the test again, except
sampling for the bits of the test pattern with some predetermined
delay for the sampling points. For instance, the delay may be,
half a clock cycle. The process continues, with the controller
5   adding another predetermined delay period on top of the previous
delay period until it detects receipt of the bit sequence of the
known test pattern. At that point, it thereafter samples received
data from that target device using the determined delay period.

Brief Description of the Drawings

Figure 1A is a timing diagram for an exemplary frame-based
digital communication protocol of the prior art.

Figure 1B is a timing diagram for a second exemplary digital
communication protocol in accordance with the prior art.

Figure 2 is a block diagram illustrating the basic
relationship of components in a computer in accordance with the
present invention.

Figure 3 is a timing diagram illustrating timing
initialization in accordance with the present invention.

20   Figure 4 is a flow chart illustrating a digital communication
initialization process in accordance with the present invention.

Detailed Description of the Invention

The invention will herein be described in connection with one or more particular embodiments adapted for use in connection with the ACR-IPB protocol. However, it should be understood that the invention has much broader applicability and can be used in connection with virtually any digital communication protocol that employs a start signal, such as a frame synchronization signal. Further, it is particularly suited for use with a digital communication system in which the start signal is sourced from a different location than the actual data that is being transmitted.

Figure 2 is a block diagram illustrating a typical architecture for a computing device utilizing the ACR-IPB communication protocol for communication with other devices. The computing device comprises a CPU 21 which communicates with a target device 25, such as a modem, through a separate controller integrated circuit 23. Typically, the controller 23 will oversee communications between the CPU and various peripheral devices. Hence, only a portion 23a of chip 23 will be dedicated to controlling ACR-IPB communications with target devices like target device 25.

In accordance with the ACR-IPB protocol, there are six lines comprising the bus 27 between the controller 23 and the target 25. The IPB bus 27 includes a transmit (TDCLK) line on which the controller places a clock signal with a maximum rate of 40 MHz. The transmit direction frame line (TDFRAME) carries the transmit

direction frame synchronization signal, which also is generated by
the controller. In one embodiment, the frame start signal
comprises a rising edge on TDFRAME. The DATAOUT line is the line
on which the actual data being transmitted from the controller to

5   the target is placed. In accordance with the ACR-IPB protocol,
separate lines are maintained for the transmit and receive
directions. Accordingly, there also are RDCLK, RDFRAME and DATAIN
lines on which corresponding signals in the receive direction,
i.e., from the target 25 to the controller 23, are carried.

10   As previously noted, the ACR-IPB protocol requires the
transmitting target device to begin driving data onto the DATAIN
line. The data must be fully valid one clock phase later in order
for the controller to read the bit correctly since the ACR-IPB
protocol specifies enter sampling (i.e., sampling on the following

15   clock edge (and assuring a 50% duty cycle). At the second clock
edge on RDCLK after the frame start and, if the first clock edge is
virtually simultaneous with the frame start signal, this provides
only slightly longer than one clock phase (e.g., half a clock
cycle) to drive the first bit of data in the frame onto DATAIN. At

20   a 40 MHz data bit rate, it can be extremely difficult for a target
device to drive the first bit of data on the DATAIN line in time to
meet this requirement. This timing problem can be alleviated if
the target device predicts the arrival of the frame start signal (a
rising edge on RDFRAME) and begins to drive data onto the DATAIN

25   line before actual receipt of the start signal. Particularly, the
start signal on RDFRAME occurs at regular intervals, and

specifically, every sixteen time slots. Accordingly, after a first
start frame signal is received by the target device on RDFRAME, the
target device can thereafter predict when the next start signal
will be received on RDFRAME. It can then set up the data for

5   transmission ahead of time so that the bit will be ready for
transmission when the second clock edge after the frame start
arrives. For instance, assuming that the target device has already
received at least one frame start signal on RDFRAME and it can
predict the arrival of future frame start signals. Accordingly, a

10   frame start signal is expected exactly sixteen clock cycles after
receipt of the last frame start signal, which we shall call time
T_0. Accordingly, at time $T_0-\gamma$, the target device begins setting up
the data for transmission. The period γ should be set as a
function of the amount of time necessary to set up the data for

15   transmission such that the target device can be ready to transmit
the data within the worst case timing scenario, namely, one half
clock cycle after time T_0.

     In at least one embodiment, the target device also can begin
transmitting the data before the second clock edge after the start

20   signal. However, it typically should not be necessary to begin
transmitting the data before the appropriate clock edge.
Specifically, the issue with respect to timely placing the first
bit of the frame on the DATAIN line is primarily a function of the
time required to initially set up the data. Once the data are set

25   up, the target device should be able to begin transmitting the
first bit (as well as all subsequent bits in the frame) at the

appropriate clock edge and that data should be valid within half a clock cycle, (when the controller samples for the data).

However, the second timing issue mentioned above is an issue for all transmitted bits. As previously noted, the transmitting

5   target device is supposed to begin transmitting the bits at the leading edges of the clock cycles and the controller will sample for the bits of the frame at the center of each clock cycle. If the round trip delay from the controller to the target device and back, 2Δ, plus the time for the data to become fully valid, δ,

10 exceeds half a clock cycle, the data will not be received at the controller at the midpoint of the clock cycle, when it is sampled. This problem is an issue with respect to every bit in the frame, but is particularly troublesome with respect to the first bit because of the extra set up time needed before the first bit can be

15 driven onto the DATAIN line.

The round trip delay is dependent on PCB implementation and internal sampling delays which likely are vendor specific. Thus, one cannot accurately characterize these delays prior to physical system test. Further, the maximum length of the board interconnect

20 between the ACR target device and the controller is not currently specified in the ACR-IPB specification. Therefore, accurate analysis of board interconnect delays is not possible. This issue is particularly important given that the ACR-IPB specification defines a maximum clock frequency of 40 MHz.

25 If the controller delays its sampling times past the mid point of the clock cycle by roughly $2\Delta+\delta$, it will sample the data

correctly for all time slots of the frame. However, as previously

noted, the delay period generally will not be known ahead of time.

Hence, in accordance with the invention, during initialization of a

communication session between the controller and a target device,

5   an appropriate delay period for sampling is determined and

thereafter applied to all transmissions from the target device and

the controller.

In one preferred embodiment of the invention, during

initialization of a communication session between the controller

10 and a target device, the controller issues an instruction to the

target device to transmit a predetermined test pattern back to the

controller in the next receive direction frame. The particular

pattern may be a full frame length or shorter. The particular

pattern can be selected in accordance with well known techniques to

15 minimize the effect of transmission errors. Both the controller

and target device know the predetermined bit pattern. The target

device, for instance, may have it permanently stored in a buffer.

Figure 3 is a timing diagram illustrating timing

initialization in accordance with the present invention upon the

20 initiation of a communication session between the controller and a

target device that will determine the round trip delay between the

controller and the target device and compensate for it.

Let us assume that the frame synchronization signal arrives at

time T_0 and the first rising clock edge on RDCLK arrives shortly

25 thereafter at time T_1. Accordingly, at the falling edge of the

clock, at time T_2, in theory, the target device should begin

transmitting the first bit of the predetermined pattern and the
controller should be able to sample correct data at time T_3.
However, in this example, let us assume that the leading edge of
the first bit does not arrive at the controller (for the various
5 delay reasons previously discussed) until shortly after time T_4.
Let us further assume that the first four bits of the predetermined
pattern are 1010 and that this is the pattern that the controller
is expecting to detect. The controller will sample the DATAIN line
at the rising edge of each clock pulse on RDCLK, namely, times T_3,
10 T_5, T_7, T_9, T_{11}, ... It can be seen in Figure 3 that the controller
will not detect the first four bits of the test pattern, 1010, but
will instead detect an indeterminate value at the first sampling
point at T_3 then will detect 101 at times T_5, T_7, T_9 and T_{11}.
Accordingly, it will determine that it has not received the correct
15 pattern 1010... but has instead received ?101....

    If the delay $2\Delta+\delta$ had been less than half a clock cycle, the
controller would have detected the correct predetermined bit
pattern and no further steps would be necessary. However, since it
did not detect the predetermined bit pattern, it will again send
20 the instruction to the target device to send the predetermined
pattern at the beginning of the next frame. In the next frame, the
controller will set its sample points with a delay. In this
example, the delay, γ, is half a clock cycle. Accordingly, instead
of sampling for the predetermined pattern with sampling points T_3,
25 T_5, T_7, T_9, T_{11},..., it will sample at points T_4, T_6, T_8, T_{10}, ... Again,
however, as can be seen in Figure 3, the controller will not detect

the correct pattern, the sampling points still being too early since the first bit of the test pattern does not arrive until shortly after time T_4. Accordingly, the controller will run the test once again, this time adding 2γ as a delay to its sampling

5   points. In this third test, the sampling points will be T_5, T_7, T_9, T_{11}, ... As can be seen in Figure 3, this time, the controller will detect the correct pattern 1010.... The controller will then know 2γ is an appropriate delay period to compensate for the delay $2\Delta+\delta$ inherent in receiving data from this particular target device.

10   The controller will thereafter utilize 2γ as a delay period for all further transmissions received from this particular target device.

    In the example of Figure 3, the sampling delay increment was half a clock cycle. However, the increment can be any amount less than one full clock cycle. Since each data bit is driven onto

15   DATAIN for one clock cycle, if the delay increment is greater than one full clock cycle, it would be possible to completely jump over the correct sampling point. Also, it is not necessary that the increment be a constant. For instance, it may be desirable for the increment to decrease with each subsequent test. Increments of one

20   half clock cycle or less should work well.

    In one preferred embodiment of the invention, each delay period is tested twice. Only if the pattern is detected both times is the corresponding delay period deemed correct. This feature can help eliminate false positives that may occur due to glitches in

25   the system, such as temporary noise or where a particular delay period puts the sampling point right on the edge of proper timing.

Since each data bit is driven an DATAIN for a full clock cycle, if the delay increments are substantially smaller than one clock cycle, there will be several sampling delay intervals that will provide accurate sampling. Accordingly, using a delay interval

5   that is greater that the absolute minimum delay interval necessary will have no adverse effect.

In other embodiments, the controller may send only one instruction to the target device to send the predetermined bit pattern per delay period, but the controller tests for receipt of

10   the pattern in the next two frames, rather than just the next frame. This embodiment is beneficial in that, in some circumstances, it is possible that the target device may take a longer time to initiate than the controller such that, the target device cannot respond to the instruction from the controller to

15   send the pre-determined pattern by the time the next frame sync signal arrives. In such a situation, the controller may not receive the predetermined pattern, not because it is sampling too early, but because the target device did not initialize fast enough to send the pattern in the next frame after the controller issued

20   the instruction, but instead sends the pattern in the next frame after that one. Thus, testing for the pattern in two connections frames can eliminate the problem.

Figure 4 is a flow diagram illustrating the basic steps performed by the controller in accordance with a basic embodiment

25   of the present invention. The process starts at step 400 upon initialization of a communication session between the controller

and a particular target device. In step 402, the controller issues
an instruction to the target device to transmit the test pattern
commencing in the first time slot of the next frame (as determined
by the next rising edge on RDFRAME). In step 404, the controller
samples for the test pattern in the next frame as dictated as
RDFRAME. In step 406, the controller compares the data detected in
step 404 to the pre-determined test pattern. If no match is
detected, flow proceeds to step 408 in which the controller adds a
pre-determined delay, γ, to the sampling points.   Flow then
proceeds back to step 402 where the test is repeated. Flow
proceeds through steps 402, 404, 406 and 408 until the pre-
determined pattern is detected in step 406. At that point, flow
proceeds to step 410, where the delay period γ used in the last
trip through step 408 is set for use in all future communications
with that target device. The process ends at step 412.

The embodiment described above in which the controller issues
an instruction for the target to send a predetermined bit pattern
is merely an exemplary embodiment. Embodiments are envisioned in
which no instruction from the controller is necessary and the
target simply sends the test patterns responsive to frame start
signals on RDFRAME. The controller can sample for the test pattern
continuously for a pre-determine number of frames before
incrementing the delay period and continuous sampling.

The appropriate delay period can and typically would be
different for each target device. Thus, the controller would need

to include logic or software for keeping track of the appropriate delay period for each target device with which it communicates.

Embodiments are envisioned in which, in order to simplify processing at the controller, after determining an appropriate delay period, the controller sends that data to the target device and the target device stores it and sends it to the controller at the beginning of each communication session as part of initialization. In this manner, each target device stores its own delay period for communicating with the controller and the controller is relieved of the burden of storing the various delay periods for all of the target devices with which it communicates.

Having thus described a few particular embodiments of the invention, various alterations, modifications, and improvements will readily occur to those skilled in the art. Such alterations, modifications and improvements as are made obvious by this disclosure are intended to be part of this description though not expressly stated herein, and are intended to be within the spirit and scope of the invention. Accordingly, the foregoing description is by way of example only, and not limiting. The invention is limited only as defined in the following claims and equivalents thereto.